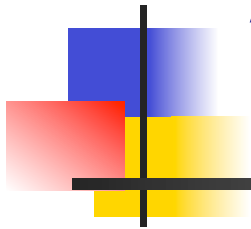


Rethinking the Rationality Postulates for Argumentation-Based Inference



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Rationality postulates (Caminada & Amgoud AIJ 2007)

- Let E be any acceptable set of arguments and

- $\text{Conc}(E) = \{\phi \mid \phi = \text{Conc}(A) \text{ for some } A \in E\}$

- Then E satisfies:

- **direct consistency** iff $\text{Conc}(E)$ does not contain two formulas ϕ and $\neg\phi$
- **strict closure** iff any ϕ deductively implied by $\text{Conc}(E)$ is in $\text{Conc}(E)$
- **indirect consistency** iff the strict closure of $\text{Conc}(E)$ is directly consistent

Assumes a deductive (i.e. strict) consequence notion



The ASPIC+ framework



- **Arguments:** DAGs where
 - Nodes are statements in some logical language \mathcal{L} containing \neg
 - Links are applications of inference rules
 - Strict rules \rightarrow
 - Defeasible rules \Rightarrow
- Constructed from consistent subsets of a **knowledge base** $\mathcal{K} \subseteq \mathcal{L}$
 - Certain premises \mathcal{K}_n + uncertain premises \mathcal{K}_p
- **Attack:**
 - On uncertain premises
 - On defeasible inferences (undercutting)
 - On conclusions of defeasible inferences (rebutting)
- **Defeat:** attack + argument ordering
- **Argument evaluation** with Dung (1995)

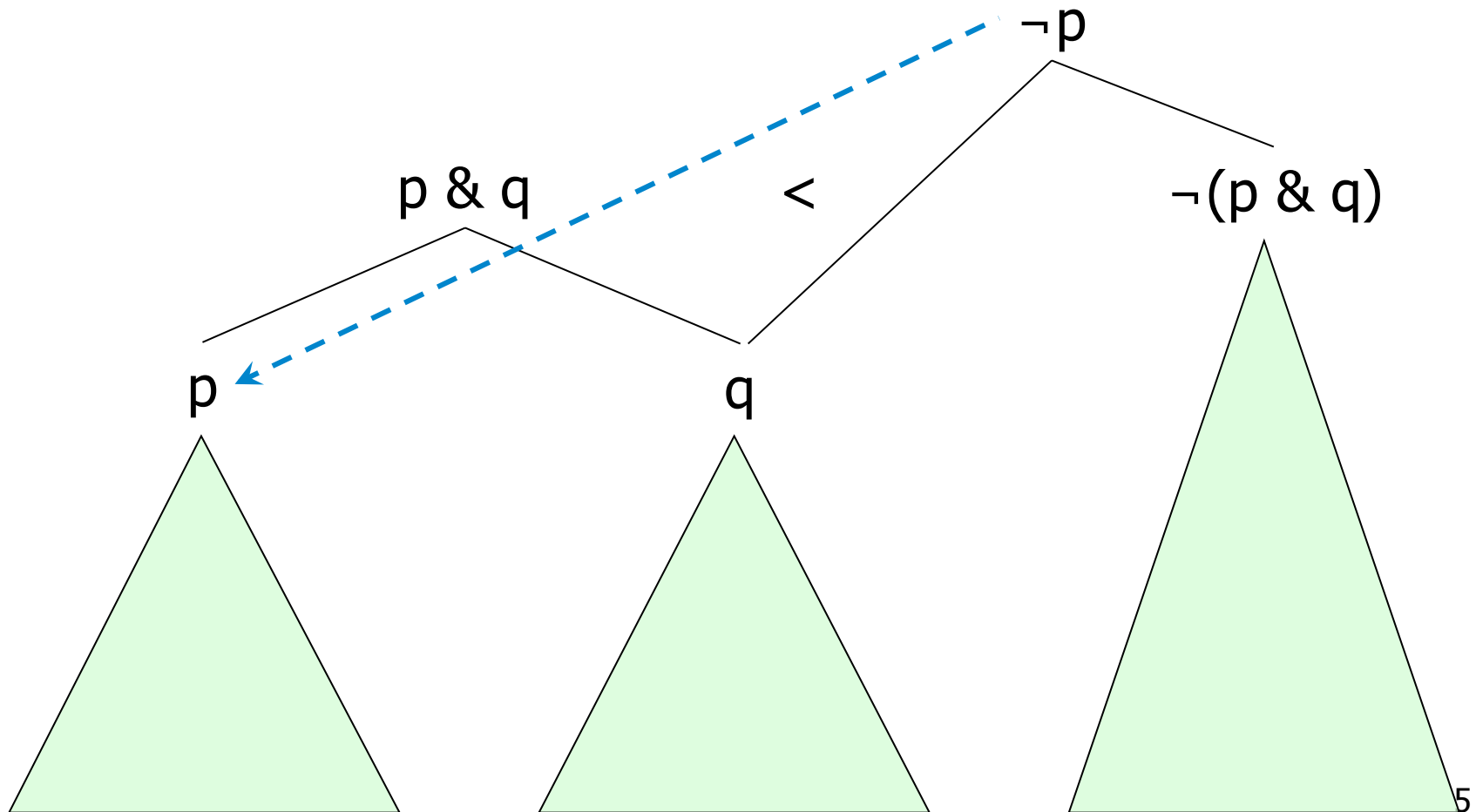
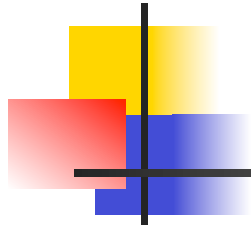
An argument is:
Fallible if it can be attacked
Infallible otherwise



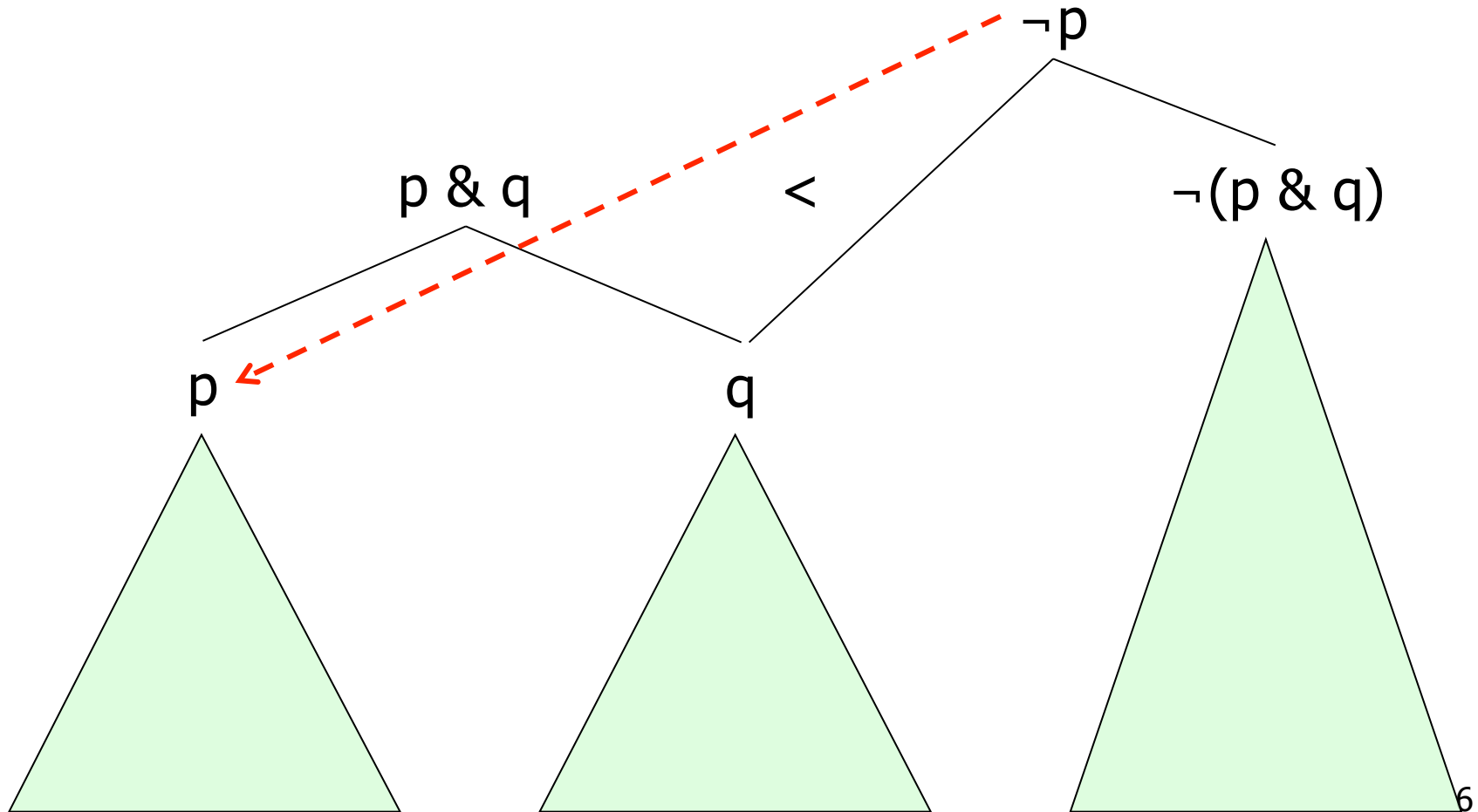
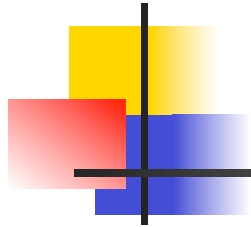
Crucial insight

- That deduction preserves truth **does not imply** that deduction preserves justification
- So that deduction preserves justification should **be independently argued**
- But deduction applied to more than one fallible subargument **can weaken** an argument,
- Since it can **aggregate** the amount of fallibility of the subarguments

contrapositive deductive reasoning



Reasonable argument ordering



The lottery paradox (Kyburg 1960)



- Assume:
 1. A lottery with 1 million tickets and 1 prize.
 2. The probability that some ticket wins is 1
 3. The probability that a given ticket T_i wins is 0.000001.
- Suppose:
 - a highly probable belief is **justified**; and
 - what can be **deduced** from a set of justified beliefs is **justified**.
- Then $\{1,2,3\}$ yields an **inconsistent** set of justified beliefs

T_1 will win and
the other tickets
will not win

The lottery paradox in ASPIC+

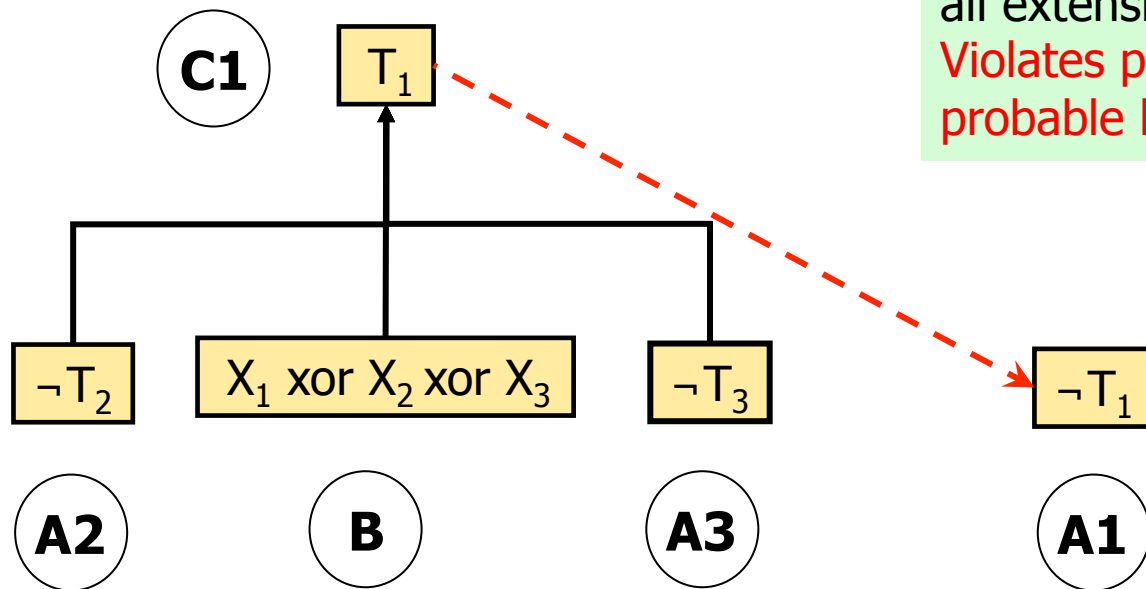
$$\begin{aligned}\mathcal{K}_p &= \{\neg T_1, \dots, \neg T_{1.000.000}\} \\ \mathcal{K}_n &= \{X_1 \text{ xor } \dots \text{ xor } X_{1.000.000}\}\end{aligned}$$

$$\begin{aligned}(\mathcal{R}_s &= \{S \rightarrow \phi \mid S \vdash_{\text{PL}} \phi \text{ and } S \text{ is finite}\} \\ \mathcal{R}_d &= \emptyset\end{aligned}$$

- **Define:** ϕ is **justified** iff some argument for ϕ is in all S-extensions

$$\mathcal{K}_p = \{\neg T_1, \neg T_2, \neg T_3\}$$

$$\mathcal{K}_n = \{X_1 \text{ xor } X_2 \text{ xor } X_3\}$$



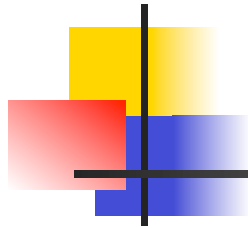
Option 1:

$C1 \approx A1$

But then for all i : $C_i \approx A_i$

So none of $\{A1, A2, A3\}$ are in all extensions

Violates principle that highly probable beliefs are justified



Excluded by
third condition
on <

$$\begin{aligned} \mathcal{K}_3 &= \{\neg T_1, \neg T_2, \neg T_3\} \\ &= \{X_1 \text{ xor } X_2 \text{ xor } X_3\} \end{aligned}$$

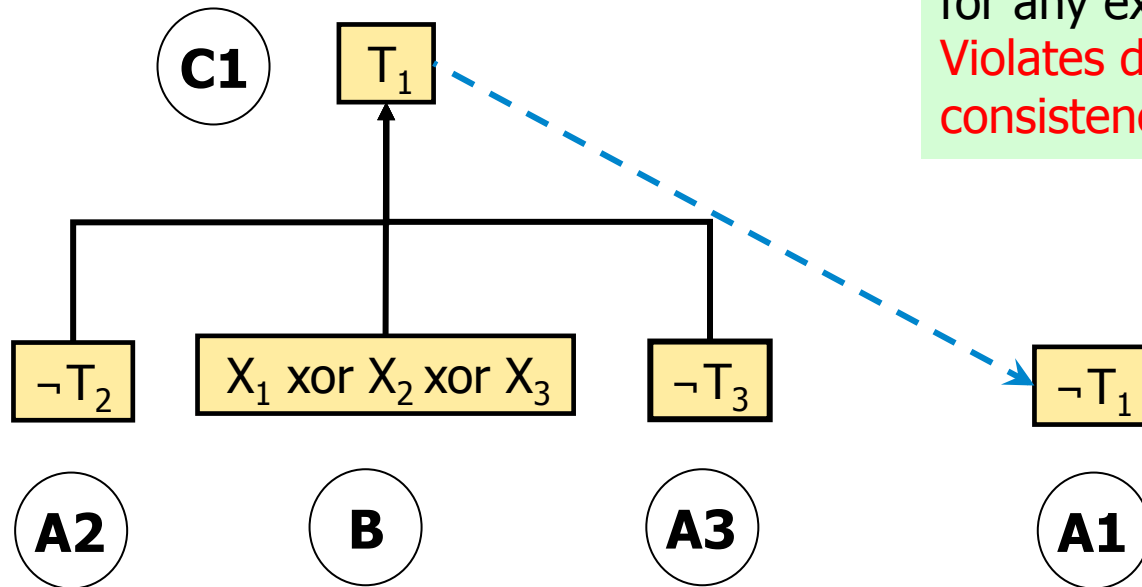
Option 2:

$C1 < A1$

But then for all i : $C_i < A_i$

So $\{A1, A2, A3, B, C1, C2, C3\} \subseteq E$
for any extension E

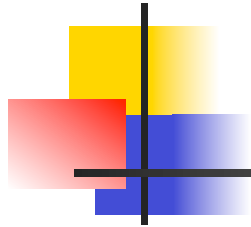
Violates direct and indirect
consistency



A decorative graphic on the left side of the slide, consisting of overlapping yellow, red, and blue squares with a black crosshair.

New rationality postulates

- Direct consistency should still hold
- Strict closure and indirect consistency should be restricted to any $S \subseteq E$ with at most one fallible argument.



Changes in ASPIC+

- Allow rebuttal on any strict inference applied to **at least two** fallible arguments
- Drop third condition on $<$

Theorem:

If strict reasoning contraposes,
and for any argument A , $\text{Premises}(A) \cup \mathcal{K}_n$ is indirectly
consistent
and conditions (1) and (2) on $<$ are satisfied

Then

direct consistency, restricted strict closure and restricted
indirect consistency are satisfied

$$\mathcal{K}_p = \{\neg T_1, \neg T_2, \neg T_3\}$$

$$\mathcal{K}_n = \{X_1 \text{ xor } X_2 \text{ xor } X_3\}$$

Option 2 again:

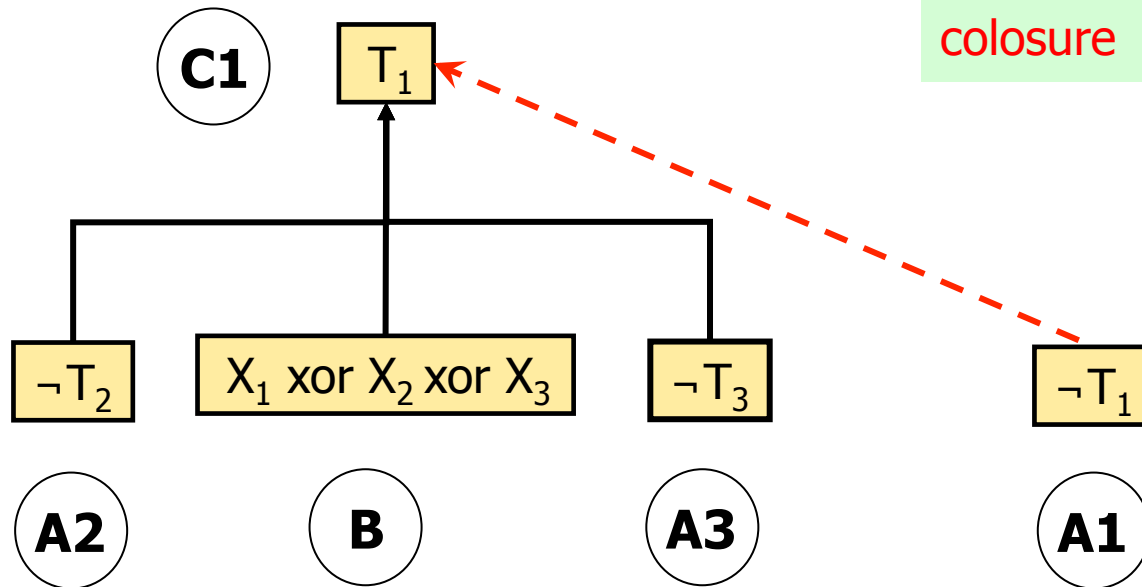
$C_i < A_i$

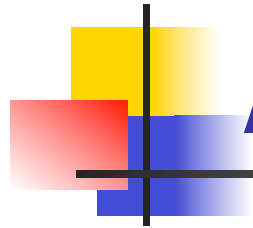
Then for all i : $C_i < A_i$

So A_1, A_2, A_3 and B are in extension E , but C_1, C_2 and C_3 are not

Violates indirect but not direct consistency

Satisfies restricted strict colasure





Added value of argumentation

- Deduction is still available in argument construction
 - Applications without attackers are still justified
 - Cannot be undercut
 - applications to a limited number of fallible subarguments can be justified, depending on the argument ordering